

Simple homogeneous structures

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Scandinavian Logic Symposium 2014
25-27 August in Tampere

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I will present some results in the intersection of these areas, i.e. we consider structures that are **both simple and homogeneous**. References (containing more references) follow at the end.

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- ① \mathcal{M} has elimination of quantifiers.
- ② Every isomorphism between finite substructures of \mathcal{M} can be extended to an automorphism of \mathcal{M} .
- ③ \mathcal{M} is the Fraïssé limit of an *amalgamation class*.

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Via the Engeler, Ryll-Nardzewski, Svenonious characterization of ω -categorical theories:

every infinite homogeneous structure has ω -categorical complete theory.

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The following classes of structures, to mention a few, have been *classified*, where 'homogeneous' implies 'countable', and 'countable' includes 'finite':

- 1 homogeneous partial orders (Schmerl 1979).
- 2 homogeneous (undirected) graphs (Gardiner, Golfand – Klin, Sheehan, Lachlan – Woodrow 1974–1980).
- 3 homogeneous tournaments (Lachlan 1984).
- 4 homogeneous directed graphs (Cherlin 1998).
- 5 homogeneous stable V -structures for any finite relational vocabulary V (Lachlan, Cherlin... 80ies).
- 6 homogeneous multipartite graphs (Jenkinson, Truss, Seidel 2012).

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Note: The case 4 contains *uncountably* many structures, by a well-known result of Henson (1972).

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- T is **supersimple** \iff the **SU-rank is ordinal valued** for every type of T , and
- T is **1-based** \iff the notion of **(in)dependence** behaves “**nicely**” on all models of T .
- T has **trivial dependence** if whenever $\mathcal{M} \models T$, $A, B, C \subseteq \mathcal{M}^{\text{eq}}$ (\mathcal{M} extended by imaginaries) and A is dependent on B over C , then there is $b \in B$ such that A is dependent on $\{b\}$ over C .

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An **infinite structure** is ω -**categorical**, **simple**, **supersimple**, **1-based** or **has trivial dependence** if its complete theory has the corresponding property.

Simple theories/structures (continued)

The SU-rank of a structure is the supremum (if it exists) of the SU-ranks of all 1-types of its complete theory.

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Two facts:

- 1 If T is ω -categorical and supersimple with finite SU-rank, then T is 1-based if and only if every definable (with parameters) $A \subseteq M^{\text{eq}}$ with SU-rank 1 is 1-based for any choice of $\mathcal{M} \models T$.
- 2 If \mathcal{M} is homogeneous, simple and 1-based, then it is supersimple with finite SU-rank and has trivial dependence.

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In fact we have:

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This implies that knowledge about properties such as 1-basedness and trivial dependence for binary simple homogeneous structures can be derived from the corresponding properties of definable sets of SU-rank 1.

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\mathcal{M} is a **binary random structure** if it does **not** have a minimal forbidden configuration of cardinality ≥ 3 .

Example: random graph.

Canonically embedded structures

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Suppose that $A \subseteq M$ and that $C \subseteq M^{\text{eq}}$ is A -definable (i.e. definable with parameters from A).

The **canonically embedded structure of \mathcal{M}^{eq} over A with universe C** is the structure \mathcal{C} which for every $0 < n < \omega$ and A -definable relation $R \subseteq C^n$ has a relation symbol which is interpreted as R (and \mathcal{C} has no other symbols).

Note that for all $0 < n < \omega$ and all $R \subseteq C^n$,
 R is \emptyset -definable in $\mathcal{C} \iff R$ is A -definable in \mathcal{M}^{eq} .

Simple homogeneous structures and canonically embedded structures with rank 1

Let \mathcal{M} and \mathcal{N} be two structures which need **not** necessarily have the same vocabulary.

\mathcal{N} is a **reduct** of \mathcal{M} if $M = N$ and for all $0 < n < \omega$ and all $R \subseteq M^n$: R is \emptyset -definable in $\mathcal{N} \implies R$ is \emptyset -definable in \mathcal{M} .

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Theorem 2 [3] *Suppose that \mathcal{M} is a binary, homogeneous, simple structure with trivial dependence. Let $A \subseteq M$ be finite and suppose that $C \subseteq M^{\text{eq}}$ is A -definable and only finitely many 1-types over \emptyset are realized in C . Assume that $\text{SU}(c/A) = 1$ for every $c \in C$, where $\text{SU}(a/A)$ is the SU-rank of the type $\text{tp}(a/A)$. Let \mathcal{C} be the canonically embedded structure of \mathcal{M}^{eq} over A with universe C . Then \mathcal{C} is a reduct of a binary random structure.*

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Note: If \mathcal{M} is homogeneous, simple and 1-based, then dependence is trivial. All known homogeneous simple structures are 1-based.

Homogeneous, simple and 1-based structures

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\mathcal{M} is **primitive** if there is no nontrivial equivalence relation on M which is \emptyset -definable.

Fact, which is straightforward to prove: *Suppose that \mathcal{M} is homogeneous (and simple) and has a nontrivial equivalence relation $E \subseteq M^2$. Let N be one of the E -classes. Then the substructure of \mathcal{M} with universe N is homogeneous (and simple).*

Homogeneous, simple and 1-based structures (cont.)

A. Aranda Lopez [1] has shown:

If \mathcal{M} is a binary, primitive, homogeneous, simple and 1-based structure with SU-rank 1, then \mathcal{M} is a binary random structure.

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A. Aranda Lopez [1] has shown:

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If we remove the assumption about the SU-rank being 1 we get:

Theorem 3. [4] *Suppose that \mathcal{M} is a binary, primitive, homogeneous, simple and 1-based structure. Then \mathcal{M} is **strongly interpretable in a binary random structure.***

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That \mathcal{M} is **strongly interpretable in \mathcal{N}** roughly means that there are integers k_1, \dots, k_m such that every element $a \in M$ can be identified with a k_i -tuple $\bar{b}_a \in N^i$ for some i in such a way that each \emptyset -definable relation in \mathcal{M} can be identified with an \emptyset -definable relation on tuples from N corresponding to elements in \mathcal{M} .

The well-known notion of **interpretability** is a generalization of strong interpretability.

[1] Andrés Aranda López, *Omega-categorical Simple Theories*, Ph.D. thesis, University of Leeds, 2014.

The following (submitted) articles can be found via the link <http://www2.math.uu.se/~vera/research/index.html> and on *arXiv*.

[2] V. Koponen, *Binary simple homogeneous structures are supersimple with finite rank*.

[3] Ove Ahlman, V. Koponen, *On sets with rank one in simple homogeneous structures*.

[4] V. Koponen, *Homogeneous 1-based structures and interpretability in random structures*.

More references are found in the sources above.